
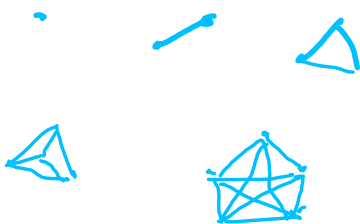
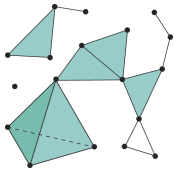



The goal of today's lecture is to give a shortcut for computing the chromatic polynomial of many graphs.

Geometry	Graph theory
<p>simplex (triangle, tetrahedron, etc.)</p> 	<p>Complete graph</p> 
<p>simplicial complex (gluing simplices together) Every face of a simplicial complex is a simplex.</p> 	<p>Chordal graph, i.e. graph that admit a simplicial elimination ordering.</p> 

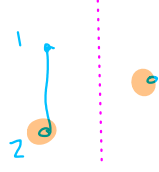
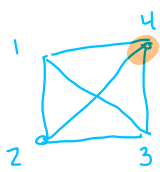
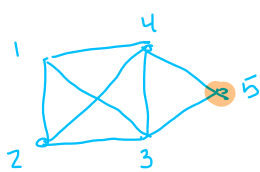
A vertex is simplicial if its neighbors form a clique.



● Simplicial

● Not simplicial

A simplicial elimination ordering is an ordering of the vertices  $v_n, \dots, v_1$  such that the vertex  $v_i$  is simplicial in  $G - \{v_n, \dots, v_{i+1}\}$



## Example

②

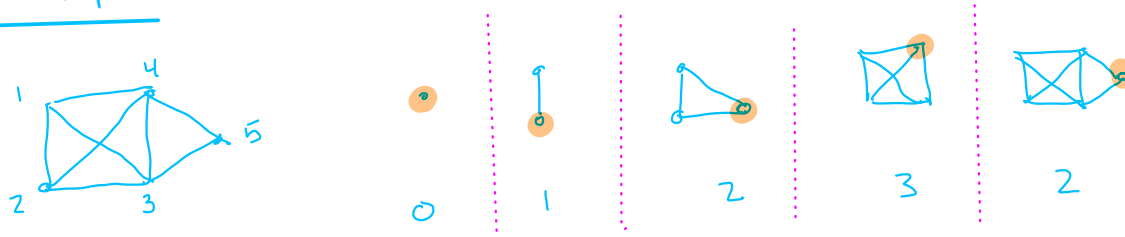
- For complete graphs, any ordering is a simplicial elimination ordering.
- For trees, an ordering starting at the leaves and going inwards is a simplicial elimination ordering.
- The cycle of length 4 has no simplicial ordering.



## Proposition

The chromatic polynomial of a graph with simplicial ordering  $v_1, \dots, v_n$  is the product of the  $k - d^*(v_i)$ 's, where  $d^*(v_i)$  is the degree of  $v_i$  in the induced subgraph with vertices  $\{v_1, \dots, v_i\}$ .

## Example



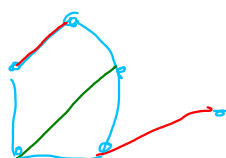
$$\chi(G; k) = k(k-1)(k-2)^2(k-3)$$

## Sketch of proof

When it is added,  $v_i$  has  $d^*(v_i)$  neighbors, all colored with different colors, so there are  $k - d^*(v_i)$  options for coloring it. Doing this process, we count all the options for coloring the graph with  $k$  colors.

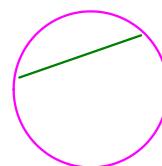
## Chordal graphs

A chord of a cycle  $C$  is an edge not in  $C$  whose endpoints are in  $C$ .



Chord

Not a chord



A cycle is chordless if it has length at least 4 and no chord.

A graph is chordal if it is simple and it has no chordless cycle (as induced subgraphs).

Theorem (Dirac, 1961)

A simple graph has a simplicial elimination ordering if and only if it is a chordal graph.

Lemma (Voloshin, 1982, or Farber-Jamison, 1986)

Every chordal graph has a simplicial vertex.

The proof of the lemma is omitted. It is in the textbook as Lemma 5.3.16.

Proof of the theorem

$\Rightarrow$  (contrapositive) If it is not chordal, there is a chordless cycle  $C$ .

In  $C$ , that has length at least 4, none of the vertices is simplicial, and no vertex can be the first one to be picked in  $C$  for the ordering.

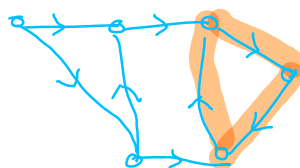
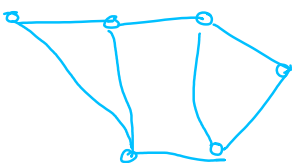
$\Leftarrow$  Using the lemma, we know that every chordal graph  $G$  has a simplicial vertex  $v$ . Delete  $v$ . Then  $G - \{v\}$  is chordal, so we can apply the lemma again, creating an ordering.



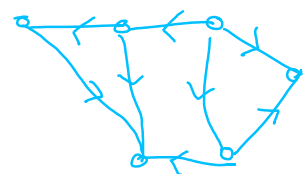
Acyclic orientations

What is the meaning of  $\chi(G; -1)$ ?

An orientation of a graph  $G$  is a digraph that has  $G$  as an underlying graph.



orientation



acyclic orientation

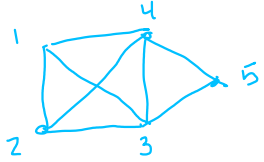
An orientation is acyclic if it has no cycle.

Theorem (Stanley, 1972)

The absolute value of  $\chi(G; -1)$  is the number of acyclic orientations of  $G$ .

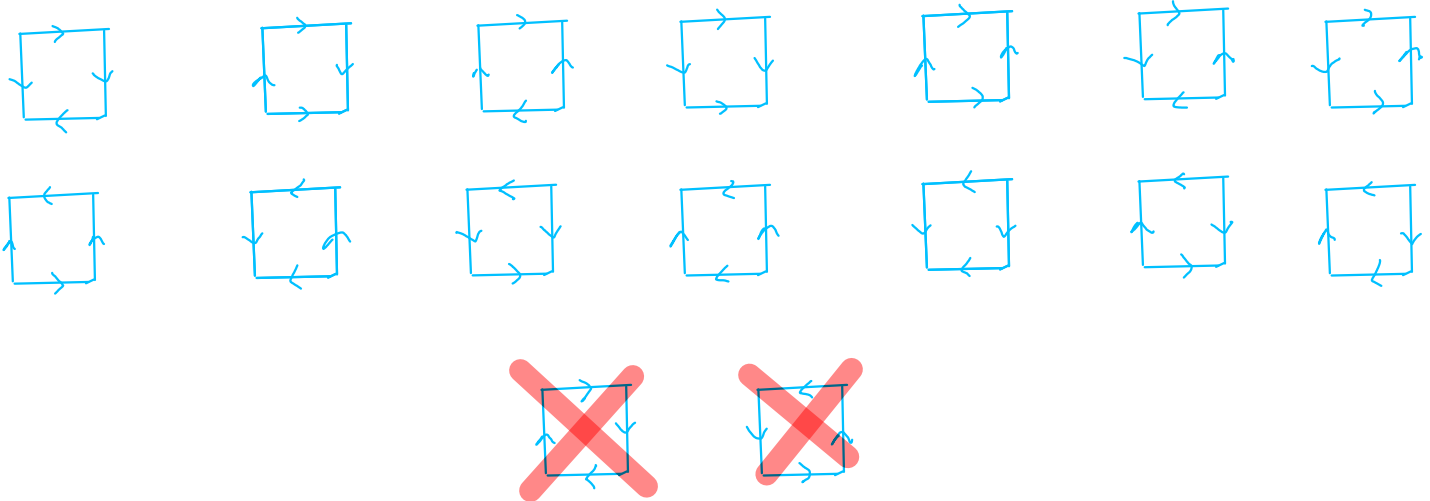
Proof is in the textbook, as the proof of Theorem 5.3.27.

### Examples



has chromatic polynomial  $\chi(G; k) = k(k-1)(k-2)^2(k-3)$ , so it has 72 acyclic orientations.

$C_4$  has chromatic polynomial  $\chi(C_4; k) = k(k-1)(k^2 - 3k + 3)$ , and  $\chi(C_4; -1)$  is 14.



Reference: Douglas B. West. Introduction to graph theory, 2nd edition, 2001.  
Section 5.3.